Efficient Data Dissemination through a Storageless Web Database

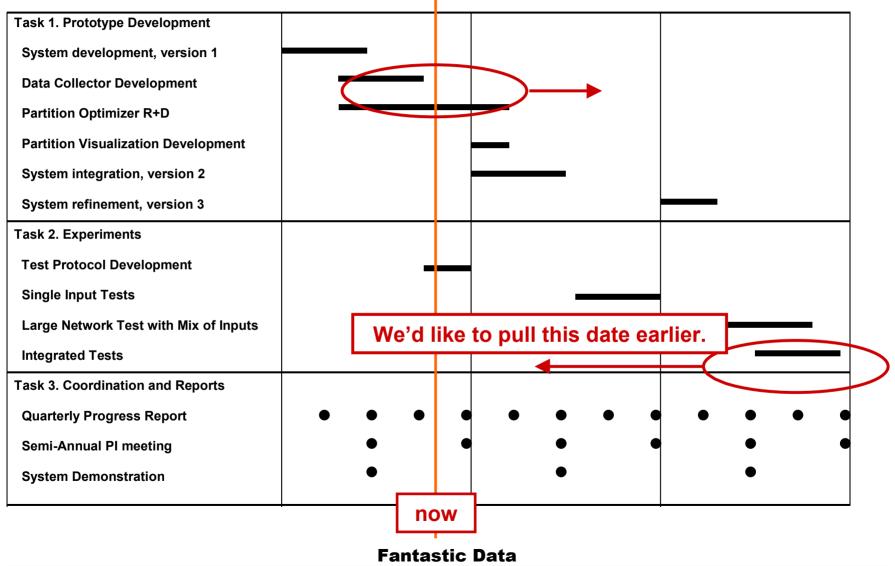
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- Automatically establish redundant data caches throughout the network based on
 - data usage patterns, transactions and queries
 - optimize cost function based on power consumption, latency, and survivability
 - no permanent storage
- Disseminate data and maintain redundant caches
 - reliable delivery on top of an unreliable channel
 - retries mitigated by
 - data expiration
 - obsolescence detection
 - priority
 - supports dynamic filter changes
 - cooperative repair

Schedule

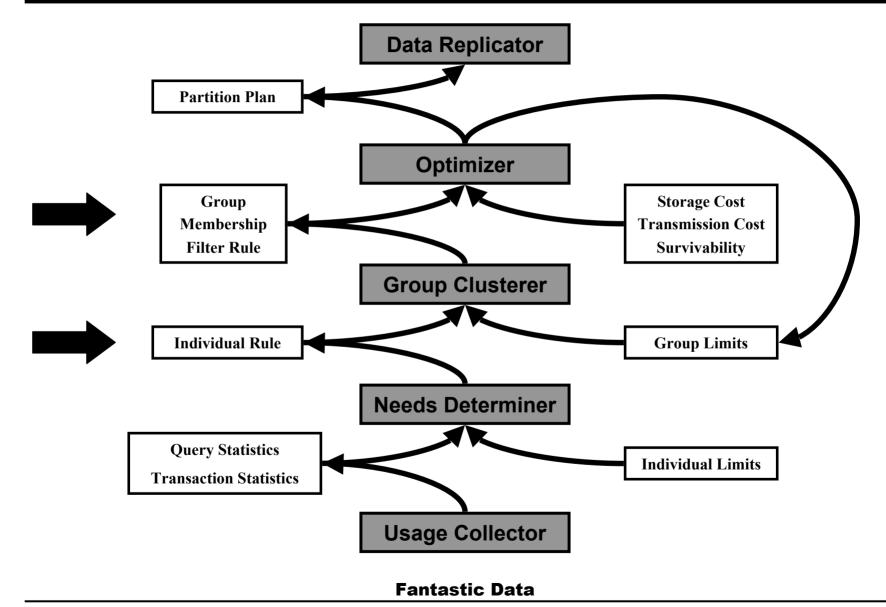
1 2 3 4 5 6 7 8 9 10 11 12 1 2 3 4 5 6 7 8 9 10 11 12 1 2 3 4 5 6 7 8 9 10 11 12



- Preparing to support experiments on the SensIT nodes
- Implementation of the in-memory cache
 - higher speed, smaller code size
 - allows easy creation of special data types, analysis functions
 - better integration of consistency bookkeeping fields with regular data fields
 - implemented distributed table creation
- Establish API
 - reduce complexity
 - eliminate features we don't really want to support
 - streamline the socket interface

- Persistent queries from application programs are used to determine dissemination filters
 - easy to "or" each application filter together
 - want to reduce composite condition complexity
 - dissemination filters are evaluated on every record change
 - reducing complexity is good, maybe even if it increases the data load
 - (a>=25 & a<=45) | (a>=30 & a<=50) is clearly equivalent to (a>=25 and a<=50), but what about
 - (power>=3.0 & lat>=39.345678 & lon>=-120.342780 & lat<=39.346001 & lon<=-120.342678 & _t>=900000000.0) | (power>=2.7 & lat>=39.345670 & lon>=-120.342778 & lat<=39.346000 & lon<=-120.342676)
 - clustering code has been developed, needs to be integrated
- Need to factor in 1-time queries
 - how often are they done?
 - how closely do they match the persistent queries?
 - how large is the remote load required to satisfy the query?

Partition Determination

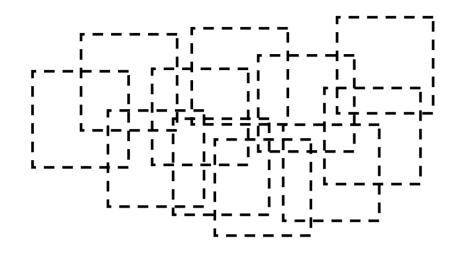


- Finish version 1
 - add single precision float data type
 - add some functions in where clause evaluator
 - min, max, mean, abs, sin, cos, tan, atan, exp, log, sqrt, ...
 - descriptive error codes
 - testing and packaging
- Need to select a distribution method that can be easily used by others and can simulate need for multihop
 - suggest broadcast to different UDP ports
 - each server broadcasts to a particular port
 - each server reads from whichever set of ports represent its neighbors
- Deliver to BBN and BAE about June 2001
 - Wider distribution later

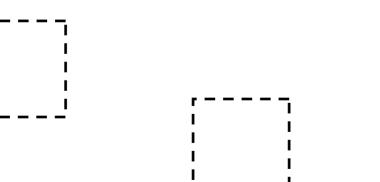
- Support use of system
 - either in operational demo, or
 - in field/lab demos as appropriate
- Integrate dynamic filter support into production version
 - code exists in larger, slower disk-resident version
- Improve filtering performance

- Support additional users
- Implement configuration options
 - data criticality
 - latency requirements
 - excess data holdback (don't need it more frequently than ...)
- Investigate relationship of Fantastic Data caches to ISI routing
 - Should we place a filtering module inside the routing layer?
 - What are the similarities/differences between our filtering approach and ISI's.

2 potentially different filtering problems

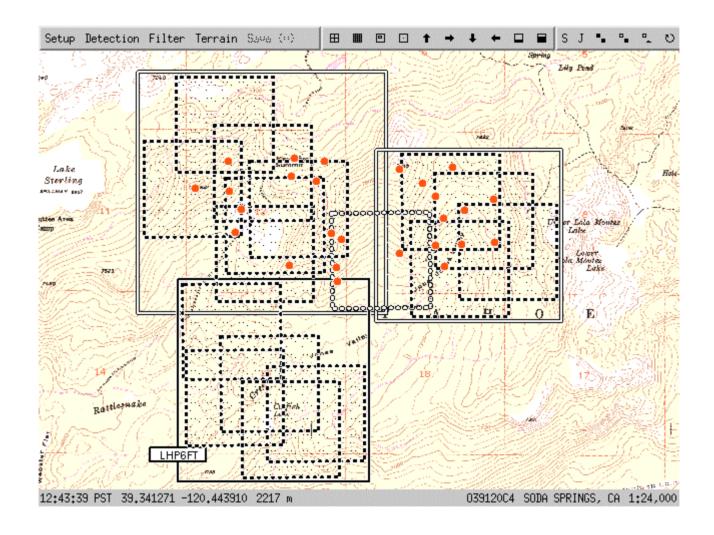


- Dense, connected interests
- Well served by local broadcast
 - flooding can be mitigated by knowledge of link state
- Data dissemination decision made by knowledge of neighbors' interest
 - can be approximated by own
- Is this the results formation problem?



- Sparse, disjoint interests
- Routing required
- Data moves across network through many uninterested nodes
- Is this the results extraction problem?

Clustering



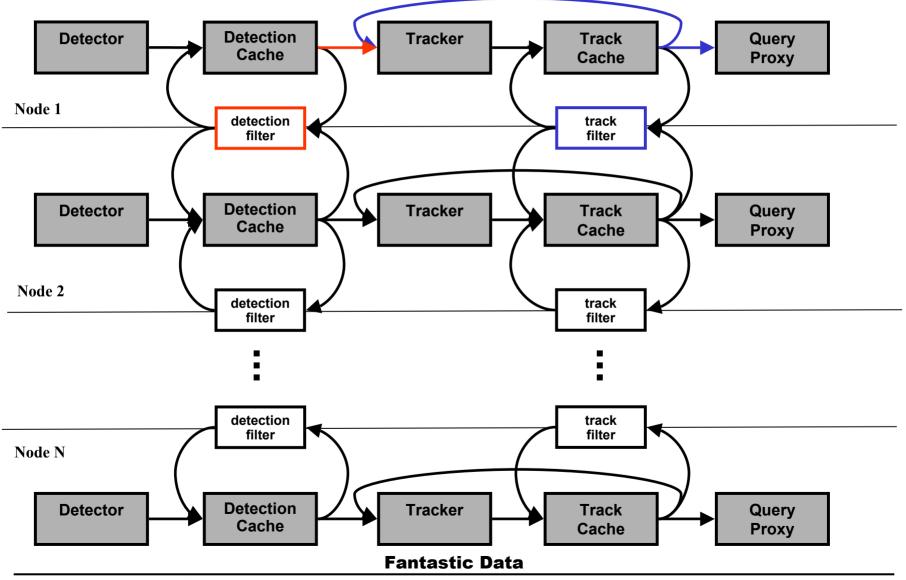
- locally determined
 - not globally optimized
 - minimize interaction between nodes required to setup filters
- incremental
 - try to disturb existing situation as little as possible
- filter tolerance
 - a little too big, a little too small, that's ok
- maintain cluster quality information
 - mean coverage of individual needs (percent, record count, bandwidth)
 - excess coverage (percent, record count, bandwidth)
 - number of members in group
 - mean age of member's input data (seconds)

- Individual rules
 - specific value, integer or character
 - mode='acoustic'
 - type='tank'
 - code=347
 - range of values (allow 1 side to be unbounded), integer or float
 - value>=3 and value<=14
 - snr>=3.5
 - power>=0.0 and power<=10.0
 - area, integer or float
 - latitude>=39.342893 and latitude<=39.358214 and longitude>=-120.451740 and longitude<=-120.430266
- Combination rules

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- "and" and "or"
 - mode='acoustic' and snr>3.5 and (latitude>=39.342893 and latitude<=39.358214 and longitude>=-120.451740 and longitude<=-120.430266)

Detection and Tracking Example



- Communicate over your own socket
 - open connection to TCP localhost:TBD or local domain /tmp/.fdb
 - send requests
 - [id] [statement];
 - where [id] is a unique integer and [statement] is any valid statement
 - new line characters may be embedded in the statement
 - it must end with a semi-colon followed by a new line character
 - cache responds with
 - [id] [status] [data]
- Use our C library functions
 - library creates and manages the socket communications
 - library provides callback functions when data arrives
 - library provides an event loop or you can use your own
 - call library function when there is activity on the socket

- Create the tables you need
 - another application or another node may have already done it
 - but doing it twice does no harm
 - unless you are first and you get it wrong
- Perform queries for configuration information
 - for example, to find out the capabilities of the node
- Perform watch operations to monitor the data you need
 - data is provided as it changes
 - you get the data before the other application gets confirmation
 - watch operations are the primary determinant of distribution filters
- Sit back and relax
 - data will come to you
- Wake up
 - process the incoming data
 - create some new data of your own

```
/*
 * Connect to the server. This is automatically done by CachePerform()
 * if you don't do it first. If a function is provided it is called
 * whenever the connection status changes.
 *
 * If socket>=0, the connection is up and this is the socket number.
 * You may use this in select() or poll().
 * If socket<0, the connection is down, and this is an error code.
 *
 * If the connection should fail, it is automatically restored and
 * any operations in progress are automatically resubmitted unless
 * CacheDisconnect() is called.
 */
extern int CacheConnect(void (*function)(int socket));
/*
 * Disconnect from the server. If a function was provided to CacheConnect(),
 * it is called when the connnection status changes.
```

```
*/
```

extern void CacheDisconnect();

API: performing an operation

```
/*
 * Ask to have an operation performed. Automatically connects
 * to the server if that has not already been done.
 *
 * Returns code assigned to this operation. Code is a positive integer.
 * A negative return indicates an error.
 *
 * Callback function is called once for each row (status>0), error (status<0),
 * and when done (status==0).
 */
extern int CachePerform(char *statement,
    void (*function)(void *ptr, void *arg), void *arg);</pre>
```

API: Statements that maintain tables

/*	
*	Statements may be any of the following:
*	CREATE TABLE table (field type, , field type, PRIMARY(field list));
*	Create a new table. The primary key is required.
*	DROP TABLE table;
*	Delete a table and all of its contents.
	DESCRIBE TABLE table;
*	Return the specification of a table.
*,	

API: Statements that change data

/*	
*	Statements may be any of the following:
*	INSERT INTO table field list VALUES value list;
*	Insert a row into the table. The row must not exist.
	PUT INTO table field list VALUES value list;
*	Insert a row into the table. If the row already exists, this
*	operation is treated as the equivalent update.
*	UPDATE table SET field list = value list WHERE condition;
*	Update all existing rows of the table that meet the condition.
*	DELETE FROM table WHERE condition;
*	Delete all existing rows that meet the condition. Caution: if no
*	condition is specified, deletes all rows of the table.
	UNDELETE FROM table WHERE condition;
*	Restore all deleted rows that meet the condition.
	PURGE FROM table WHERE condition;
*	Permanently removes all rows of the table that meet the condition.
*	Caution: The presence of deleted data is necessary to ensure
*	consistency of the redundant caches. Data should be purged only
*	if some other mechanism makes sure that all nodes purge the same
*	data at the same time. For example, a data expiration feature
*	based on time of day and time of data creation is safe.
*,	

API: Statements that query data

/*	
/	
*	Statements may be any of the following:
*	SELECT field list FROM table WHERE condition;
*	Return the rows of the table that match the condition.
	WATCH field list FROM table WHERE condition;
*	Return the rows of the table that match the condition as they are
*	inserted, updated, or deleted.
	SELECT AND WATCH field list FROM table WHERE condition;
*	Do both select and watch operations. Guaranteed to return all of
*	the existing rows before any changes and to not miss any changes.
	CANCEL operation;
*	Cancel a previously specified operation, for example, an
*	ongoing watch operation.
*	
,	

API: Interpreting the results

```
/*
 * Returns the operation code associated with the given result pointer.
 * Returns a negative error code if something is wrong.
 */
extern int CacheResultCode(void *ptr);
/*
 * Returns the status associated with the given result pointer.
 * See enum CacheStatus. 0 means the operation is done.
 * A positive status means that there is valid data.
 * A negative return is an error.
 */
extern int CacheResultStatus(void *ptr);
/*
 * Return the number of rows changed or returned by the operation.
 * A negative return is an error code.
 */
extern int CacheResultImpact(void *ptr);
```

API: Interpreting the results

```
/*
 * Returns a pointer to the data associated with the result pointer.
 * If the result status is positive, this is the row data in text form.
 * If the result status is negative, this is an error message.
 * Returns 0 if there is something wrong.
 */
extern char *CacheResultMessage(void *ptr);
/*
 * Returns the number of fields in the result. Valid only if the
 * associated operation was a guery.
 *
 * A negative return is an error code.
 */
extern int CacheResultNvalue(void *ptr);
/*
 * Returns the value of the specified field in the specified result
 * in text form.
 * Returns 0 if there is something wrong.
 */
extern char *CacheResultValue(void *ptr, int it);
```

API: data field types

```
/*
  Data types include:
 *
      integer (32 bits),
 *
      short (16 bits),
 *
      byte (8 bits),
 *
      char (variable length),
 *
      double (64 bits),
 *
      blob (variable length, uninterpreted binary data).
 *
 */
```

API: where clauses

/*
 * Conditions may be composed of any field, constants, and the operators
 *
 * +, -, *, /, &, |, !, ^, * <, <=, =, >=, >, !=.
 *
 * The absence of a WHERE clause is interpreted as whatever portion of
 * the table is maintained on the local node. This is not the same as
 * all data in the table throughout the system.
 *
 * Where clauses on queries (especially watch queries) are the major
 * determinant of data distribution filters.
 */

API: special data fields

/*	
*	All tables are automatically supplied with several special fields. The
*	values of these fields may not be set with insert, update, or put
*	statements, but they may be returned by queries and used in conditions.
*	The special fields are the following:
*	_b integer, the node that changed the record.
*	_s integer, the series of the record change.
*	_n integer, the sequence number of the record change.
*	_t double, the time at which the record change was originated.
*	_pb integer, the node that previously changed the record.
*	_ps integer, the series of the previous record change.
*	_pn integer, the sequence number of the previous record change.
*	_pt double, the time at which the previous record change was originated.
*	_lt double, the time at which the record change was performed on
*	the local node.
*,	

Location of nodes

create table node (id integer, latitude double, longitude double, primary(id));

Capabilities of nodes

create table capability (id integer, type char, ..., primary(id));

Detections

create table detection (id integer, latitude double, longitude double, start integer, end integer, cpa integer, power double, primary(id));

Tracks

create table track (id integer, type char, latitude double, longitude double, confidence double, primary(id));